N-Queens Combinatorial Problem Polyglot FP for Fun and Profit – Haskell and Scala 🔉 📕

first see the problem solved using the List monad and a Scala for comprehension

then see the Scala program translated into Haskell, both using a do expressions and using a List comprehension

understand how the Scala for comprehension is desugared, and what role the withFilter function plays

also understand how the Haskell do expressions and List comprehension are desugared, and what role the guard function plays

Part 1





Miran Lipovača



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Soo

Programming in

Fourth Edition



A comprehensive step-by-step guide

Martin Odersky Lex Spoon Bill Venners

artima









Yes, I know, Lex Spoon looks a bit too young in that photo, and Bill Venner's hair doesn't look realistic in that drawing. Martin Odersky's drawing is not bad at all.

2 @philip_schwarz

By the way, don't be put off by the fact that the book image below is of the fourth edition. Excerpts from the **N-queens** section of that edition are great, and by the way, in the fifth edition (new for **Scala 3**) I don't see that section (it may have moved to upcoming book **Advanced Programming in Scala**). Also, I do use **Scala 3** in this slide deck.









Martin Odersky

Bill Venners







A particularly suitable application area of **for expressions** are **combinatorial puzzles**.

An example of such a puzzle is the 8-queens problem: Given a standard chess-board, place eight queens such that no queen is in check from any other (a queen can check another piece if they are on the same column, row, or diagonal).

Here on the right is a sample solution to the puzzle.





To find a solution to this problem, it's actually simpler to generalize it to chessboards of arbitrary size.

Hence the problem is to place N queens on a chess-board of $N \ge N$ squares, where the size N is arbitrary.

We'll start numbering cells at one, so the upper-left cell of an $N \ge N$ board has coordinate (1,1) and the lower-right cell has coordinate (N,N).

To solve the N-queens problem, note that you need to place a queen in each row. So you could place queens in successive rows, each time checking that a newly placed queen is not in check from any other queens that have already been placed.

In the course of this search, it might happen that a queen that needs to be placed in row k would be in check in all fields of that row from queens in row 1 to k - 1. In that case, you need to abort that part of the search in order to continue with a different configuration of queens in columns 1 to k - 1.



And below is an example in which there is nowhere to place queen number 6 because every other cell on the board is in **check** from queens 1 to 5.





An **imperative solution** to this problem would place queens one by one, moving them around on the board.

But it looks difficult to come up with a scheme that really tries all possibilities.

<u>A more functional approach represents a solution directly, as a value</u>. <u>A solution consists of a list of coordinates, one for each queen placed on the board</u>.

Note, however, that a full solution can not be found in a single step. It needs to be built up gradually, by occupying successive rows with queens.

This suggests a recursive algorithm.

Assume you have already generated all solutions of placing k queens on a board of size $N \ge N$, where k is less than N.

Each solution can be represented by a list of length k of coordinates (row, column), where both row and column numbers range from 1 to N.

It's convenient to <u>treat these partial solution lists as stacks</u>, where the coordinates of the queen in row k come first in the list, followed by the coordinates of the queen in row k - 1, and so on.

The bottom of the stack is the coordinate of the queen placed in the first row of the board.

All solutions together are represented as a list of lists, with one element for each solution.



See below for the list representing the above solution to the **8-queens** puzzle.



List((8, 6), (7, 3), (6, 7), (5, 2), (4, 8), (3, 5), (2, 1), (1, 4))





Now, to place the next queen in row k + 1, generate all possible extensions of each previous solution by one more queen.

This yields another list of solutions lists, this time of length k + 1.

Continue the process until you have obtained all solutions of the size of the chess-board N.

This **algorithmic idea** is embodied in function **placeQueens** below:







The outer function **queens** in the program above simply calls **placeQueens** with the size of the board **n** as its argument.

The task of the function application placeQueens(k) is to generate all partial solutions of length k in a list.

Every element of the list is one solution, represented by a list of length k. So placeQueens returns a list of lists.



List((8, 6), (7, 3), (6, 7), (5, 2), (4, 8), (3, 5), (2, 1), (1, 4)) If the parameter k to placeQueens is 0, this means that it needs to generate all solutions of placing zero queens on zero rows. There is only one solution: place no queen at all. This solution is represented by the empty list. So if k is zero, placeQueens returns List(List()), a list consisting of a single element that is the empty list.

Note that this is quite different from the empty list List(). If placeQueens returns List(), this means no solutions, instead of a single solution consisting of no placed queens.





In the other case, where k is not zero, all the work of **placeQueens** is done in the **for expression**.

The **first generator** of that **for expression** iterates through all solutions of placing k - 1 queens on the board.

The second generator iterates through all possible columns on which the k^{th} queen might be placed.

The **third part** of the **for expression** defines the newly considered queen position to be the pair consisting of row k and each produced column.

The **fourth part** of the **for expression** is a filter which checks with **isSafe** whether the new queen is safe from check by all previous queens (the definition of **isSafe** will be discussed a bit later.)

If the new queen is not **in check** from any other queens, it can form part of a partial solution, so **placeQueens** generates with queen::queens a new solution.

If the new queen is not safe from **check**, the filter returns **false**, so no solution is generated.

Here is the **for expression** again, for reference.



for {
 queens <- placeQueens(k - 1)
 column <- 1 to n
 queen = (k, column)
 if isSafe(queen, queens)
} yield queen :: queens</pre>

The only remaining bit is the **isSafe** method, which is used to check whether a given queen is **in check** from any other element in a list of queens.

Here is the definition:





The **isSafe** method expresses that a queen is safe with respect to some other queens if it is not **in check** from any other queen.

The **inCheck** method expresses that queens q1 and q2 are mutually **in check**.

It returns true in one of three cases:

- 1. If the two queens have the same row coordinate.
- 2. If the two queens have the same column coordinate.
- 3. If the two queens are on **the same diagonal** (*i.e.*, the difference between their rows and the difference between their columns are the same).

The first case – that the two queens have the same row coordinate – cannot happen in the application because **placeQueens** already takes care to place each queen in a different row. So you could remove the test without changing the functionality of the program.





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Algorithm

We can solve this problem with a recursive algorithm:

- Suppose that we have already generated all the solutions consisting of placing k-1 queens on a board of size n.
- Each solution is represented by a list (of length k-1) containing the numbers of columns (between 0 and n-1).
- The column number of the queen in the k-1th row comes first in the list, followed by the column number of the queen in row k-2, etc.
- The solution set is thus represented as a set of lists, with one element for each solution.
- Now, to place the kth queen, we generate all possible extensions of each solution preceded by a new queen:

Example: N-Queens

The eight queens problem is to place eight queens on a chessboard so that no queen is threatened by another.

In other words, there can't be two queens in the same row, column, or diagonal.



list (0, 3, 1)

List (2,0,2,1)

We now develop a solution for a chessboard of any size, not just 8.

One way to solve the problem is to place a queen on each row.

Once we have placed $k\,$ - 1 queens, one must place the kth queen in a column where it's not "in check" with any other queen on the board.



List(5, 2, 6, 1, 7, 4, 0, 3)

In online course **FP** in Scala, the board coordinates are **zero-based**, and since it is obvious that the nth queen is in row n, a solution is a list of **column indices** rather than a list of **coordinate pairs**.



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Because a column index is sufficient to indicate the coordinates of the nth queen, we can see that the type of **isSafe**'s first parameter is now **Int**, i.e. a column index, rather than (**Int**, **Int**) i.e. a coordinate pair.

Also, the type of a solution is List[Int] rather than List[(Int, Int)], and solutions are returned in a Set, rather than a List, to distinguish between a single solution, which is an ordered list, and the collection of all solutions, which is not ordered (though it could be, if we so wanted).



Here are the two versions of the program side by side.



```
def queens(n: Int): List[List[(Int,Int)]] = {
    def placeQueens(k: Int): List[List[(Int,Int)]] =
        if (k == 0)
        List(List())
    else
        for {
           queens <- placeQueens(k - 1)
           column <- 1 to n
           queen = (k, column)
           if isSafe(queen, queens)
        } yield queen :: queens
    placeQueens(n)
    }
</pre>
```

```
def isSafe(queen: (Int, Int), queens: List[(Int, Int)]) =
   queens forall (q => !inCheck(queen, q))

def inCheck(q1: (Int, Int), q2: (Int, Int)) =
   q1._1 == q2._1 || // same row
   q1._2 == q2._2 || // same column
   (q1._1 - q2._1).abs == (q1._2 - q2._2).abs // on diagonal
```



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```
def isSafe(col: Int, queens: List[Int]): Boolean = {
  val row = queens.length
  val queensWithRow = (row - 1 to 0 by -1) zip queens
  queensWithRow forall
    { case (r, c) => col != c && math.abs(col - c) != row - r }
}
```





Same as on the previous slide, but with really minor tweaks thanks to Scala 3



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```
def queens(n: Int): List[List[(Int,Int)]] =
  def placeQueens(k: Int): List[List[(Int,Int)]] =
   if k == 0
   then List(List())
    else
      for
        queens <- placeQueens(k - 1)</pre>
        column < -1 to n
        queen = (k, \text{ column})
        if isSafe(queen, queens)
      yield queen :: queens
  placeQueens(n)
def isSafe(queen: (Int, Int), queens: List[(Int, Int)]) =
 queens forall (q => !inCheck(queen, q))
def inCheck(queen1: (Int, Int), queen2: (Int, Int)) =
 queen1(0) == queen2(0) || // same row
 queen1(1) == queen2(1) || // same column
  (queen1(0) - queen2(0)).abs == (queen1(1) - queen2(1)).abs // on diagonal
```

```
def queens(n: Int): Set[List[Int]] =
 def placeQueens(k: Int): Set[List[Int]] =
   if k == 0
    then Set(List())
    else
      for
        queens <- placeQueens(k - 1)</pre>
        col <- ⊘ until n
        if isSafe(col, queens)
      yield col :: queens
  placeQueens(n)
def isSafe(col: Int, queens: List[Int]): Boolean =
 val row = queens.length
 val queensWithRow = (row - 1 to 0 by -1) zip queens
  queensWithRow forall
    { case (r, c) => col != c && math.abs(col - c) != row - r }
```

In the next part of this slide deck, we are going to be using the program on the right.

The **queens** and **placeQueens** functions are pretty much the same as in the second version on the previous slide, but board coordinates are one-based, and the two functions return a **List** rather than a **Set**.

The logic for determining if a queen is safe from previously placed queens, is organised in a way that is in part a hybrid between the two programs on the previous slide.



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```
def queens(n: Int): List[List[Int]] =
  def placeQueens(k: Int): List[List[Int]] =
    if k == 0
    then List(List())
    else
       for
        queens <- placeQueens(k - 1)
        queen <- 1 to n
        if safe(queen, queens)
        yield queen :: queens</pre>
```

placeQueens(n)



List(6, 3, 7, 2, 8, 5, 1, 4)

```
def safe(queen: Int, queens: List[Int]): Boolean =
  val (row, column) = (queens.length, queen)
  val safe: ((Int,Int)) => Boolean = (nextRow, nextColumn) =>
    column != nextColumn && !onDiagonal(column, row, nextColumn, nextRow)
  zipWithRows(queens) forall safe
```

```
def onDiagonal(row: Int, column: Int, otherRow: Int, otherColumn: Int) =
    math.abs(row - otherRow) == math.abs(column - otherColumn)
```

```
def zipWithRows(queens: List[Int]): Iterable[(Int,Int)] =
  val rowCount = queens.length
  val rowNumbers = rowCount - 1 to 0 by -1
  rowNumbers zip queens
```

```
assert(queens(4) == List(List(3, 1, 4, 2), List(2, 4, 1, 3)))
assert(queens(6) == List(List(3, 6, 2, 5, 1, 4),
                    List(4, 1, 5, 2, 6, 3),
                    List(5, 3, 1, 6, 4, 2),
                    List(2, 4, 6, 1, 3, 5)))
```



The first thing we are going to do, is translate our **N-queens** program from **Scala** into **Haskell**.

Let's start with the **queens** and **placeQueens** functions.



In **Scala** we use a **for comprehension**, whereas in **Haskell** we use a **do expression**.

In Scala, the filtering is done by

if safe(queen, queens)

whereas in **Haskell** it is done by

guard (safe queen queens)



If you are asking yourself what the **guard** function is, don't worry: we'll be looking at it soon.



We can simplify the **Haskell** version a little bit by eliminating the calls to **guard** and **return**, which can be achieved by switching from a **do expression** to a **list comprehension**.





```
queens n = placeQueens n
where
placeQueens 0 = [[]]
placeQueens k =
   [queen:queens |
   queens <- placeQueens(k-1),
   queen <- [1..n],
   safe queen queens]</pre>
```



If you are asking yourself how the **list comprehension** eliminates the need for **guard** and **return**, don't worry: we'll be looking at that soon.



And here is the translation of the whole program

```
def queens(n: Int): List[List[Int]] =
  def placeQueens(k: Int): List[List[Int]] =
    if k == 0
    then List(List())
    else
        for
        queens <- placeQueens(k - 1)
        queen <- 1 to n
        if safe(queen, queens)
        yield queen :: queens
    placeQueens(n)</pre>
```

```
queens n = placeQueens n
where
placeQueens 0 = [[]]
placeQueens k = [queen:queens |
        queens <- placeQueens(k-1),
        queen <- [1..n],
        safe queen queens]</pre>
```

def safe(queen: Int, queens: List[Int]): Boolean =
 val (row, column) = (queens.length, queen)
 val safe: ((Int,Int)) => Boolean = (nextRow, nextColumn) =>
 column != nextColumn && !onDiagonal(column, row, nextColumn, nextRow)
 zipWithRows(queens) forall safe

def onDiagonal(row: Int, column: Int, otherRow: Int, otherColumn: Int) =
 math.abs(row - otherRow) == math.abs(column - otherColumn)

```
def zipWithRows(queens: List[Int]): Iterable[(Int,Int)] =
  val rowCount = queens.length
  val rowNumbers = rowCount - 1 to 0 by -1
  rowNumbers zip queens
```

safe queen queens = all safe (zipWithRows queens)
where
 safe (r,c) = c /= col && not (onDiagonal col row c r)
 row = length queens
 col = queen

onDiagonal row column otherRow otherColumn =
 abs (row - otherRow) == abs (column - otherColumn)

```
zipWithRows queens = zip rowNumbers queens
where
rowCount = length queens
rowNumbers = [rowCount-1,rowCount-2..0]
```



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```
scala> queens(4)
val res0: List[List[Int]] = List(List(3, 1, 4, 2), List(2, 4, 1, 3))
scala> queens(5)
val res1: List[List[Int]] = List(List(4, 2, 5, 3, 1), List(3, 5, 2, 4, 1), List(5, 3, 1, 4, 2), List(4, 1, 3, 5, 2),
            List(5, 2, 4, 1, 3), List(1, 4, 2, 5, 3), List(2, 5, 3, 1, 4), List(1, 3, 5, 2, 4),
            List(3, 1, 4, 2, 5), List(2, 4, 1, 3, 5))
scala> queens(6)
val res2: List[List[Int]] = List(List(5, 3, 1, 6, 4, 2), List(4, 1, 5, 2, 6, 3),
            List(3, 6, 2, 5, 1, 4), List(2, 4, 6, 1, 3, 5))
```

haskell> queens 4
[[3,1,4,2],[2,4,1,3]]
haskell> queens 5
[[4,2,5,3,1],[3,5,2,4,1],[5,3,1,4,2],[4,1,3,5,2],
[5,2,4,1,3],[1,4,2,5,3],[2,5,3,1,4],[1,3,5,2,4],
[3,1,4,2,5],[2,4,1,3,5]]
haskell> queens 6

[[5,3,1,6,4,2],[4,1,5,2,6,3], [3,6,2,5,1,4],[2,4,6,1,3,5]]





For N > 6, the number of solutions is too large for the solutions to be displayed simply by printing them to the screen.

Let's at least check that we get the expected number of solutions.

N	N-queens solution count	15	2,279,184	<pre>scala> queens(8).size val res0: Int = 92 scala> queens(9).size val res1: Int = 352</pre>	Haskell>	length	(queens	<mark>8</mark>)
1	1	16	14,772,512		92			
2	0	17	95.815.104		Haskell>	length	(queens	<mark>9</mark>)
3	0				352			
4	2	18	666,090,624	scala queens (10) size	Haskell>	length	(queens	10)
5	10	19	4,968,057,848	val res2: Int = 724	724	rengen	queens	10)
6	4	20	39,029,188,884	<pre>scala> queens(11).size val res3: Int = 2680</pre>	Haskell> length (que 2680	length	(queens	11)
7	40	21	314,666,222,712					
8	92	22	2,691,008,701,644	<pre>scala> queens(12).size</pre>	Haskell>	length	(queens	12)
9	352	23	24.233.937.684.440	val res4: Int = 14200	14200			
10	724			s_{cala} queens (13) size	Hackella	longth	(augons	13)
11	2,680	24	227,514,171,973,736	val res5: Int = 73712	73712	Length	(queens	13)
12	14 200	25	2,207,893,435,808,352					
13	73,712	26	22,317,699,616,364,044	<pre>scala> queens(14).size val res6: Int = 365596</pre>	Haskell> 365596	length	(queens	14)
14	365,596	27	234,907,967,154,122,528					



In the next five slides, we are going to look at how the Scala for comprehension is implemented under the hood.

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Here is how the **N-queens for comprehension**, without the filter, is implemented, i.e. what it looks like when it is **desugared** (when its syntactic sugar is removed).





placeQueens(k - 1) flatMap { queens =>
 (1 to n) map { queen =>
 queen :: queens
}



To aid our understanding, here is a similar **for comprehension** that is simpler to understand, but can be run at the **REPL**, together with its **desugared** version.

for

nums <- List(List(1,2),List(3,4))
num <- List(10,20)
yield num :: nums</pre>





And here is the value produced by both the **for comprehension** and its **desugared** equivalent.

List(List(10, 1, 2), List(20, 1, 2), List(10, 3, 4), List(20, 3, 4))



Same as on the previous slide, except that now we are adding a filter to the **for comprehension**.

When the filter is **desugared**, it results in a call to **List's withFilter** function.





```
List(List(20, 1, 2), List(20, 3, 4))
```

Let's see what the withFilter function does:

- it is applied to a monad, in this case a List
- it takes a **predicate**, i.e. a function that takes a parameter and returns either **true** or **false**.
- it applies the **predicate** to each value yielded by the **monad** (in this case, each value contained in a **List**)
- it filters out the values for which the **predicate** returns **false**
- It lets through the values for which the **predicate** returns **true**

On the next slide, we look at an example of how the behaviour of a **for comprehension** (or its **desugared** equivalent) is affected by the behaviour of the **withFilter** function.





List(List(10, 1, 2), List(20, 1, 2), List(10, 3, 4), List(20, 3, 4))

withFilter lets through no elements



List()

Notice that the **predicate** function passed to **withFilter** may or may not use its parameter. Like the **map** and **flatMap** functions, **withFilter** is used to bind the value(s) yielded by a **monad** (if any) to a variable name which is then available for use in subsequent computations.



The difference is that while the functions passed to **map** and **flatMap** are almost certain to use the variable, in the case of **withFilter**, it may choose to influence the overall computation without actually using the variable.

List(List(1,2),List(3,4)) flatMap { nums =>

List(10,20) withFilter { num =>

false

} map { num =>

num :: nums



Before we move on, here is a final, illustration of what the withFilter function does.

```
scala> List(1,2,3).withFilter(n => true).map(identity)
val res0: List[Int] = List(1, 2, 3)
```

```
scala> List(1,2,3).withFilter(n => n > 1).map(identity)
val res1: List[Int] = List(2, 3)
```

```
scala> List(1,2,3).withFilter(n => false).map(identity)
val res2: List[Int] = List()
```



Earlier we saw that the **do expression** that is the **Haskell** translation of the **Scala for comprehension**, involves the use of the **guard** function.

Also, later we intend to look at how the **Haskell do expression** is implemented under the hood, which involves both the **guard** function and the (>>) function.

So, in the next three slides, we are going to look at how Miran Lipovača explains guard and (>>).

Learn You a Haskell for Great Good!

A Beginner's Guide

Miran Lipovača

Here is the default implementation of the (>>) function:

(>>) :: (Monad m) => ma -> mb -> mb m >> n = m >>= \setminus -> n >>= is Haskell's equivalent of Scala's
flatMap and _ -> n is Haskell's
equivalent of Scala's anonymous
function _=> n.



Normally, passing some value to a function that ignores its parameter and always returns some predetermined value always results in that predetermined value. With **monads** however, their context and meaning must be considered as well.

Here is how >> acts with Maybe: Here on the **ghci**> [] >> [3] ghci> Nothing >> Just 3 right is how >> [] Nothing acts with lists. **ghci**> [3] **>>** [4] ghci> Just 3 >> Just 4 [4] Just 4 ghci> [3] >> [] ghci> Just 3 >> Nothing Nothing scala> None >> 3.some scala> List() >> List(3) val res0: Option[Int] = None val res0: List[Int] = List() scala> 3.some >> 4.some scala> List(3) >> List(4) val res1: Option[Int] = Some(4) val res1: List[Int] = List(4) scala> 3.some >> None scala> List(3) >> List() val res2: Option[Nothing] = None val res2: List[Nothing] = List() scala> None >> None scala> List() >> List() val res3: Option[Nothing] = None val res3: List[Nothing] = List()

Miran Lipovača

import cats._
import cats.implicits._

Same as above, but using

the **Scala Cats** library:

MonadPlus and the guard Function

List comprehensions allow us to filter our output. For instance, we can filter a list of numbers to search only for numbers whose digit contains a 7:

```
ghci> [x | x <- [1..50], '7' `elem` show x]
[7,17,27,37,47]</pre>
```

We apply **show** to **x** to turn our number into a string, and then we check if the character '7' is part of that string. To see how filtering in list comprehensions translates to the list **monad**, we need to check out the **guard** function and the **MonadPlus** type class. The **MonadPlus** type class is for **monads** that can also act as **monoids**. Here is its definition:

```
class Monad m => MonadPlus m where
  mzero :: ma
  mplus :: ma -> ma -> ma
```

mzero is synonymous with **mempty** from the **Monoid** type class, and **mplus** corresponds to **mappend**. Because lists are **monoids** as well as **monads**, they can be made an instance of this type class:

```
instance MonadPlus [] where
mzero = []
mplus = (++)
```

For lists, mzero represents a nondeterministic computation that has no results at all - a failed computation. mplus joins two nondeterministic values into one. The guard function is defined like this:

```
guard :: Bool -> m ()
guard True = return ()
guard False = mzero
```

guard takes a **Boolean** value. If that value is **True**, guard takes a () and puts it in a minimal default context that succeeds. If the **Boolean** value is **False**, guard makes a failed monadic value.

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```
Here it is in action:
```

```
ghci> guard (5 > 2) :: Maybe ()
Just ()
ghci> guard (1 > 2) :: Maybe ()
Nothing
ghci> guard (5 > 2) :: [()]
[()]
ghci> guard (1 > 2) :: [()]
[]
```

ghci> [x | x <- [1..50], '7' `elem` show x] [7,17,27,37,47]

same result

This looks interesting, but how is it useful? In the list monad, we use it to filter out nondeterministic computations:

ghci> [1..50] >>= (\x -> guard ('7' `elem` show x) >> return x)
[7,17,27,37,47]

The result here is the same as the result of our previous list comprehension. How does guard achieve this? Let's first see how guard functions in conjunction with >>:

```
ghci> guard (5 > 2) >> return "cool" :: [String]
["cool"]
ghci> guard (1 > 2) >> return "cool" :: [String]
[]
```

If guard succeeds, the result contained within it is the empty tuple. So then we use >> to ignore the empty tuple and present something else as the result. However, if guard fails, then so will the return later on, because feeding an empty list to a function with >>= always results in an empty list. guard basically says, "If this Boolean is False, then produce a failure right here. Otherwise, make a successful value that has a dummy result of () inside it." All this does is to allow the computation to continue.



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Here we see the **desugaring** of the **Scala for comprehension** again, but we also **desugar** the equivalent **Haskell do expression**.











placeQueens(k-1) >>= \queens ->
[1..n] >>= \queen ->
guard (safe queen queens) >>
return (queen:queens)



Note how, while the **desugared Scala** code introduces the use of **flatMap**, **withFilter** and **map**, the **desugared Haskell** code introduces the use of **>>=** (**Haskell**'s equivalent of **flatMap**) and **>>**.

Note also how, in **Haskell**, the **guard** function is used in both the **do expression** and in its **desugared** equivalent.



Same as on the previous slide, except that here the logic is simpler to understand and can be run at the **REPL**.

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List(List(20, 1, 2), List(20, 3, 4))



[[20,1,2],[20,3,4]]



List(List(10, 1, 2), List(20, 1, 2), List(10, 3, 4), List(20, 3, 4))



List()

```
ghci> [1,2] >>= \n -> ['a','b'] >>= \ch -> return (n,ch)
[(1,'a'),(1,'b'),(2,'a'),(2,'b')]
```

Here is the previous expression rewritten in **do notation**:

```
listOfTuples :: [(Int,Char)]
listOfTuples do
    n <- [1,2]
    ch <- ['a','b']
    return (n,ch)</pre>
```

Do Notation and List Comprehensions

Using lists with **do notation** might remind you of something you've seen before. For instance, check out the following piece of code:

ghci> [(n, ch) | n <- [1,2], ch <- ['a','b']]
[(1,'a'),(1,'b'),(2,'a'),(2,'b')]</pre>

Yes, list comprehensions! In our do notation example, n became every result from [1,2]. For every such result, ch was assigned a result from ['a', 'b'], and then the final line put (n,ch) into a default context (a singleton list) to present it as the result without introducing any additional nondeterminism.

In this **list comprehension**, the same thing happened, but we didn't need to write **return** at the end to present (**n**, **ch**) as the result, because the output part of a **list comprehension** did that for us.

In fact, list comprehensions are just syntactic sugar for using lists as monads. In the end, list comprehensions and lists in do notation translate to using >>= to do computations that feature nondeterminism.





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We saw earlier what the Haskell do expression looks like when it is desugared.

Desugaring the **list comprehension** produces the same result.





As a recap, let's see again the translation of the whole program from **Scala** to **Haskell**.

```
def queens(n: Int): List[List[Int]] =
  def placeQueens(k: Int): List[List[Int]] =
    if k == 0
    then List(List())
    else
        for
        queens <- placeQueens(k - 1)
        queen <- 1 to n
        if safe(queen, queens)
        yield queen :: queens
    placeQueens(n)</pre>
```

queens n = placeQueens n where placeQueens 0 = [[]] placeQueens k = [queen:queens] queens <- placeQueens(k-1),</pre> queen <- [1..n], safe queen queens]

def safe(queen: Int, queens: List[Int]): Boolean =
 val (row, column) = (queens.length, queen)
 val safe: ((Int,Int)) => Boolean = (nextRow, nextColumn) =>
 column != nextColumn && !onDiagonal(column, row, nextColumn, nextRow)
 zipWithRows(queens) forall safe

def onDiagonal(row: Int, column: Int, otherRow: Int, otherColumn: Int) =
 math.abs(row - otherRow) == math.abs(column - otherColumn)

```
def zipWithRows(queens: List[Int]): Iterable[(Int,Int)] =
  val rowCount = queens.length
  val rowNumbers = rowCount - 1 to 0 by -1
  rowNumbers zip queens
```

safe queen queens = all safe (zipWithRows queens)
where
 safe (r,c) = c /= col && not (onDiagonal col row c r)
 row = length queens
 col = queen

onDiagonal row column otherRow otherColumn =
 abs (row - otherRow) == abs (column - otherColumn)

```
zipWithRows queens = zip rowNumbers queens
where
rowCount = length queens
rowNumbers = [rowCount-1,rowCount-2..0]
```



To conclude the first part of this slide deck, see below for a simple way of displaying the solutions to the **N-queens** problem.



6.3 Combinatorial Search Example

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<pre>def show(queens: List[Int]) = {</pre>						
val lines =						
<pre>for (col <- queens.reverse)</pre>						
vield Vector.fill(aueens.lenath)("* ").updated(col. "X ").mkString						
"\n" + (lines mkString "\n")						
}	> show: (queens: List[Int])java.lang.Stri					
queens(4) map show	> res0: scala.collection.immutable.Set[ja					
	* * X *					
	X * * *					
	* * * X					
	* X * * ". "					
	* X * *					
	* * * X					
	Y * * *					
	× * Y * ")					



There is lots to do in part two. We'll kick off by taking the **Scala** version of the program, and extending it so that it can display a single solution board as follows:

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We'll then get the program to display, all together, the results of queens(N) for N = 4, 5, 6, 7, 8. Plus much more. See you then.